

Alethe: Towards a Generic SMT Proof Format

PxTP 2021

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Some History



Within VeriT

- ▶ First: Ad-hoc ( TACAS 2006)
- ▶ Later: Redesigned ( PxTP 2011)
- ▶ Syntax changed over time

SMTCoq

- ▶ One of the first users
- ▶ Verified checker ( CPP 2011)
- ▶ Base for automation in Coq ( CAV 2017, PxTP -3h20)

Fine-Grained Proofs and Isabelle

- ▶ Support for reasoning with bound variables ( CADE 2017, JAR 2020)
- ▶ Typical for pre-processing in SMT
- ▶ Isabelle/HOL integration ( CADE +49h40)

Now!

-  Proofonomicon
-  Speculative Specification
-  It's now Alethe! ( PxTP -0h05)

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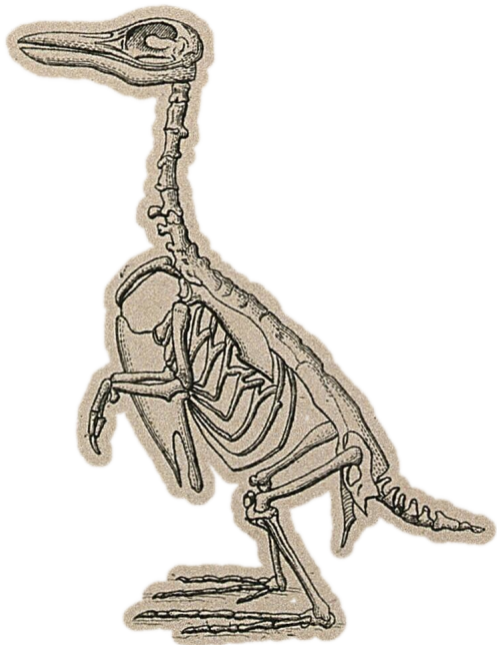
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Basic Structur

$$\frac{\begin{array}{c} t_2 \\ \hline t_3 \\ \vdots \\ t_1 \quad \neg t_1 \end{array}}{\perp} \text{resolution}$$

```
(assume a0 t1)
(assume a1 t2)
(step s1 (cl t3)
  :premises (a1)      :rule rule1)
  ...
(step s20 (cl (not t1))
  :premises (s19)     :rule rule2)
(step s21 (cl )
  :premises (a0 s20)  :rule resolution)
```


Subproofs With Assumptions

$$\frac{\frac{t_1}{t_2} \quad \frac{[t_2] \quad \vdots \quad t_3}{\neg t_2, t_3} \text{subproof}}{t_3} \text{resolution}}$$

```
(assume a0 t1)
(step s1 (cl t2)
  :premises (a0) :rule rule1)
(anchor :step s2)
  (assume s2.a1 t2)
  ...
  (step s2.s10 (cl t3)
    :premises (s2.s9) :rule rule2)
(step s2 (cl (not t2) t3) :rule subproof)
(step s3 (cl t3)
  :premises (s1 s2) :rule resolution)
```

Reasoning With Binders

$$\frac{}{x \mapsto y \triangleright x = y} \text{ refl}$$
$$\frac{}{x \mapsto y \triangleright f(x) = f(y)} \text{ cong}$$
$$\frac{}{\forall x. f(x) = \forall y. f(y)} \text{ bind}$$

```
(anchor :step s2 :args ((:= (x S) y)))  
  (step s2.s1 (cl (= x y))      :rule refl)  
  (step s2.s2 (cl (= (f x) (f y)))  
              :rule cong)  
(step s2 (cl (= (forall ((x S)) (f x))  
                (forall ((y S)) (f y))))  
        :rule bind)
```



Rules

Current State

- ▶ Overall 90 rules, mostly simple tautologies
- ▶ Seven categories – with overlaps
- ▶ Some historic overhead
- 🏗️ Cleanup and normalization

How can we accommodate different solvers?


- ▶ Some solvers might be able to use rules more strictly.
- ▶ Example:
 - ▶ $a = b \wedge b = c \rightarrow a = c$
 - ▶ $c = b \wedge a = b \rightarrow a = c$



Have an optional annotation to mark restricted usage.

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A Checker and Elaborator

- ▶ “A second pair of eyes”.
- ▶ Small, independent codebase – in Rust.
- ▶ Long term: rewrite steps to their stricter form, framework to replace non-standard rules by standard rules.



Bruno Andreotti

Support in *cvc5*

- ▶ Part of a wider effort to overhaul the proof module of *cvc5*.
- ▶ Will add more theories to Alethe.



Hanna Lachnitt and the wider *cvc5* team.

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Speculative Specification

<http://www.verit-solver.org/alethe.pdf>

Feedback

<https://gitlab.uliege.be/verit/alethe>